Analysis of the Blocking Behaviour of Schema Transformations in Relational Database Systems

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WikiMedia schema revisions:



Source: <u>http://yellowstone.cs.ucla.edu/schema-evolution/index.php/Schema_Evolution_Benchmark</u>

Curino et al.

WikiMedia schema revisions:



• 90% require a write lock.

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WikiMedia schema revisions:



- 90% require a write lock.
- Largest took 22 hours to complete for wikipedia.

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Some systems can not go offline:

- Telecom systems
- Payment systems
- Airline reservation systems
- Online services

Contributions

- An experimental investigation of the existing solutions for online schema changes.
 - What solutions are currently available?
 - To what degree do these work?
 - What are the open problems?
- A solution sketch for complex online schema transformations.

Earlier work

- Requirements for online schema changes
- Important classes of relational transformations
- A benchmark for online schema transformations based on TPC-C

Lesley Wevers, Menno Tammens, Matthijs Hofstra, Marieke Huisman and Maurice van Keulen. A Benchmark for Non-blocking Schema Transformations. DATA 2015.

Requirements for schema changes

Functionality

Transformations should:

- Provide ACID guarantees
- Be specified declaratively
- Include program updates

Performance

Transformations should:

- Have minimal impact on concurrent transactions:
 - Blocking
 - Slowdown
 - Aborts
- Commit as fast as possible

Schema Transformations

Semantic:

- Add / remove:
 - Relations
 - Columns
 - Constraints
- Change the cardinality of a relationship
- Use surrogate keys instead of natural keys

Implementation / performance:

- Add / remove indices
- Precompute aggregates
- Change normalization
- Change encoding of relationships

• ...

Benchmark

TPC-C 10 Warehouse District W*10 W History 3k 100k W*30k+ 1+ Stock Customer W*100k W*30k New-Order W*9k+ 3+ 0-1 1+ W Item Order-Line Order 100k W*30k+ W*300k+ 5-15

New order

Payment

Order status

Delivery

Stock level

Benchmark

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Benchmark Process

- Setup database
- Start TPC-C
- Intro period
- Transform:
 - Schema
 - Stored procedures
- Outro period
- Stop TPC-C



Contenders

Data definition language:

- PostgreSQL: Partial Instantaneous DDL
- MySQL: Partial Online DDL
- Oracle 11g

Ronströms method:

- Third-party tools (MySQL):
 - pt-online-schema-change
 - o oak-online-alter-table
 - Facebook's online-schema-change
- Oracle's DBMS_PARALLEL_COPY

Experimental Setup

- HammerDB + Custom driver script
- Database: 30 warehouses
- Workload: 64 threads
- Out of the box behaviour: no tuning

Experimental Results

DDL: Column Operations



DDL: Index Operations



DDL: Bulk Data Transformations



DDL: Complex Transformations



DDL Conclusions

Basic transformations:

- Mixed results
- Online DDL doesn't always do what you expect.

Complex transformations:

• Every complex operation depends on data transformations, which can not be performed online.

id	street	number
1	A str	4
2	B str	78
3	C str	5b
4	D str	1

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id	address	
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id	street	number
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2	B str	78
3	C str	5b
4	D str	1

. . .

. . .

. . .

ON UPDATE

ON INSERT

ON DELETE

id address

id	street	number
1	A str	4
2	B str	78
3	C str	5b
4	D str	1

. . .

id	address
1	A str 4
2	B str 78

- ON UPDATE ...
- ON INSERT

ON DELETE ...

id	street	number
1	A str	4
2	B str	78
3	C str	5b
4	Anotherstreet	87

id	address
1	A str 4
2	B str 78

- ON UPDATE ...
- ON INSERT ...
- ON DELETE ...

id	street	number
1	A str	4
2	B str	78
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id	address
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- ON UPDATE ...
- ON INSERT ...
- ON DELETE ...

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ON UPDATE ... ON INSERT ... ON DELETE ...

id	address
1	A str 4
2	B str 78
3	C str 5b
4	Anotherstreet 87

pt-osc: parameters



(a) Chunk size 1000, (b) Chunk size 1000, (c) Chunk size load 64 load 4 10000, load 4

pt-osc: DDL operations



Oracle's DBMS_PARALLEL_COPY



Ronstroms method: Conclusions

- Good results for basic transformations
- Support for some complex transformations: derived columns, splitting / merging tables
- Only needs support for triggers in the DBMS.
- There are no techniques for transactional composition

Conclusions

• DDL provided by DBMS has very limited support for online transformations.

• Ronström's approach is promising, but lacks support for composition.

Solution Sketch for Complex Transformations using Ronström's method

Algebraic Transformations



Propagation using Ronström's method



We need to know how to propagate updates in the source of an operator to updates in the

Propagation must be efficient: We may need to build indices before starting propagation.

Composition in Ronström's method

id	street	number	id	address		id	address
1	A str	4			a		
2	B str	78		Ť			
3	C str	5b					
4	D str	1					
ON	UPDATE .	•••					
ON	INSERT .	•••					
ON	DELETE .						

Composition in Ronström's method

id	street	number		id	addres	S		id	address
1	A str	4					_		
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ON	UPDATE .	••		ON U	JPDATE				
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ON	DELETE .	••		ON I	DELETE				

Composition in Ronström's method

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ON	UPDATE .	•••
ON	INSERT .	· · ·

ON DELETE ...

Discussion

DBMSs can provide:

- More efficient alternatives to Ronström's method.
- Instantaneous operations.
- On-the-fly operations.

Outlook

Optimizer for online transformations: generate execution plans for algebraic transformations.

- Select materialization approaches for every operator.
- Optimize for minimal time to commit.

For more information

Read the paper:

- Analysis of the Blocking Behaviour of Schema Transformations in Relational Database Systems
- Download the benchmark implementations:
- http://wwwhome.ewi.utwente.nl/~weversl2/?page=ost
- Benchmark specification:
 - Lesley Wevers, Menno Tammens, Matthijs Hofstra, Marieke Huisman and Maurice van Keulen. A Benchmark for Non-blocking Schema Transformations. DATA 2015.